

Five key results

- b  $\neg\exists\forall\phi$  is true on an interpretation I iff  $\forall\forall\neg\phi$  is true on I.

*Proof.* Suppose that  $\neg\exists\forall\phi$  is true on an interpretation I. So for every assignment  $\alpha$ ,  $\neg\exists\forall\phi$  is true on I relative to  $\alpha$ . So for every assignment  $\alpha$ ,  $\exists\forall\phi$  is false on I relative to  $\alpha$ . So for every assignment  $\alpha$ , there is no  $\forall$ -variant  $\alpha'$  of  $\alpha$  such that  $\phi$  is true on I relative to  $\alpha'$ . So for every assignment  $\alpha$ , and every  $\forall$ -variant  $\alpha'$  of  $\alpha$ ,  $\phi$  is false on I relative to  $\alpha'$ . So for every assignment  $\alpha$ , and every  $\forall$ -variant  $\alpha'$  of  $\alpha$ ,  $\neg\phi$  is true on I relative to  $\alpha'$ . So for every assignment  $\alpha$ ,  $\forall\forall\neg\phi$  is true on I relative to  $\alpha$ . So  $\forall\forall\neg\phi$  is true on I. Now suppose that  $\forall\forall\neg\phi$  is true on an interpretation I. So for every assignment  $\alpha$ ,  $\forall\forall\neg\phi$  is true on I relative to  $\alpha$ . So for every assignment  $\alpha$ , and every  $\forall$ -variant  $\alpha'$  of  $\alpha$ ,  $\neg\phi$  is true on I relative to  $\alpha'$ . So for every assignment  $\alpha$ , and every  $\forall$ -variant  $\alpha'$  of  $\alpha$ ,  $\phi$  is false on I relative to  $\alpha'$ . So for every assignment  $\alpha$ , there is no  $\forall$ -variant  $\alpha'$  of  $\alpha$  such that  $\phi$  is true on I relative to  $\alpha'$ . So for every assignment  $\alpha$ ,  $\exists\forall\phi$  is false on I relative to  $\alpha$ . So for every assignment  $\alpha$ ,  $\neg\exists\forall\phi$  is true on I relative to  $\alpha$ . So  $\neg\exists\forall\phi$  is true on I. ■

Models and consistency

- a.iii  $\{\forall x\exists y(Py \wedge \neg x = y), \exists x\exists y((Px \wedge Py) \wedge \neg x = y)\}$  is consistent.

*Proof.* Let I be the following interpretation. The domain is  $\{1, 2\}$ ; 'P' denotes  $\{1, 2\}$ . Then for everything in the domain there is some distinct thing in the domain which is in the set denoted by 'P', so  $\forall x\exists y(Py \wedge \neg x = y)$  is true on I. And there are (at least) two (distinct) things in the domain which are in the set denoted by 'P', so  $\exists x\exists y((Px \wedge Py) \wedge \neg x = y)$  is true on I. So there is an interpretation on which every wff in  $\{\forall x\exists y(Py \wedge \neg x = y), \exists x\exists y((Px \wedge Py) \wedge \neg x = y)\}$  is true, and the set is consistent. ■

- b.iii  $\{\exists xRxx, \neg a = b, \forall x\forall y(Rxy \rightarrow ((x = a \wedge y = b) \vee (x = b \wedge y = a)))\}$  is inconsistent.

*Proof.* Suppose that there is an interpretation I on which every wff in this set is true. Since  $\exists xRxx$  is true on I, there is something in the domain of I, call it Bob, such that  $\langle \text{Bob}, \text{Bob} \rangle$  is in the set denoted by 'R'. Since  $\forall x\forall y(Rxy \rightarrow ((x = a \wedge y = b) \vee (x = b \wedge y = a)))$  is true on I,  $(Rxy \rightarrow ((x = a \wedge y = b) \vee (x = b \wedge y = a)))$  is true on I no matter what is assigned to 'x' and 'y'. In particular, it is true on I when Bob is assigned to both 'x' and 'y'. Since 'Rxy' is true on I when Bob is assigned to 'x' and 'y', we must have that  $((x = a \wedge y = b) \vee (x = b \wedge y = a))$  is true on I when Bob is assigned to 'x' and 'y'. So when Bob is assigned to 'x' and 'y', either 'x = a' and 'y = b' are both true on I, or 'x = b' and 'y = a' are both true on I. But 'x = a' and 'y = b' are both true on I when Bob is assigned to 'x' and 'y' only if 'a' and 'b' both denote Bob on I, and similarly 'x = b' and 'y = a' are both true on I when Bob is assigned to

'x' and 'y' only if 'a' and 'b' both denote Bob on I. So we must have that 'a' and 'b' both denote Bob on I. But then ' $\neg a = b$ ' is false on I, which contradicts the assumption that it is true on I. So there is no such interpretation, and the set is inconsistent. ■

Extending the tableau system to QC

- g. Done on the board
- h. Done on the board